### Mask Set Errata for Mask 0N57U

This report applies to mask 0N57U for these products:

• S32K144

Erratum ID	Erratum Title
e10658	ADC: Trigger loss is not reported
e6939	Core: Interrupted loads to SP can cause erroneous behavior
e9004	Core: ITM can deadlock when global timestamping is enabled
e9005	Core: Store immediate overlapping exception return operation might vector to incorrect interrupt
e6940	Core: VDIV or VSQRT instructions might not complete correctly when very short ISRs are used
e11543	FlexCAN: Nominal Phase SJW incorrectly applied at CRC Delimiter
e10856	FTM: Safe state is not removed from channel outputs after fault condition ends if SWOCTRL is being used to control the pin
e10752	LPI2C: Slave Busy Flag may incorrectly read as zero when 10-bit address mode enabled in slave mode.
e10792	LPI2C: Slave Transmit Data Flag may incorrectly read as one when TXCFG is zero.
e11097	LPSPI: Command word not loaded correctly when TXMSK=1
e11089	LPSPI: In Continuous transfer mode with CPHA =1, WCF bit is not set for every word.
e10655	LPSPI: Module Busy Flag may incorrectly read as zero in Master mode
e10527	LPUART: Setting and immediately clearing SBK bit can result in transmission of two break characters
e10716	RTC: Timer Alarm Flag can assert erroneously
e10777	SCG: Corrupted status when the system clock is switching.
e11063	SMC: An asynchronous wakeup event during VLPS mode entry may result in possible system hang scenario.
e11114	SMC: invalid data might be fetched while accessing Flash in VLP modes

### Table 1. Errata and Information Summary



Table 2.	Revision	History
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Revision	Changes
06/Mar/2019	Initial revision

### e10658: ADC: Trigger loss is not reported

**Description:** A trigger loss will not be reported in the following cases:

1. The time between two consecutive triggers on a trigger input is less than 4 cycles of ADC's Module interface clock.

2. One or multiple triggers come on any input while a trigger either stored or under process for that input in the trigger handler.

3. The trigger on a particular input asserts during the period of end-of-conversion cycle of a conversion initiated from this input.

- In these cases the Error in Multiplexed Trigger Request (ADC\_SC2[TRGSTERR]) bit in the ADC Status and Control Register 2 will not be set.
- For the third case the Trigger status (ADC\_SC2[TRGSTLAT]) bit in the ADC Status and Control Register 2 will also not be set.
- **Workaround:** The time between two consecutive triggers on a particular input should be kept more than one ADC conversion time.

### e6939: Core: Interrupted loads to SP can cause erroneous behavior

Description: Arm Errata 752770: Interrupted loads to SP can cause erroneous behavior

This issue is more prevalent for user code written to manipulate the stack. Most compilers will not be affected by this, but please confirm this with your compiler vendor. MQX<sup>™</sup> and FreeRTOS<sup>™</sup> are not affected by this issue.

Affects: Cortex-M4, Cortex-M4F

Fault Type: Programmer Category B

Fault Status: Present in: r0p0, r0p1 Open.

If an interrupt occurs during the data-phase of a single word load to the stack-pointer (SP/ R13), erroneous behavior can occur. In all cases, returning from the interrupt will result in the load instruction being executed an additional time. For all instructions performing an update to the base register, the base register will be erroneously updated on each execution, resulting in the stack-pointer being loaded from an incorrect memory location.

The affected instructions that can result in the load transaction being repeated are:

- 1) LDR SP,[Rn],#imm
- 2) LDR SP,[Rn,#imm]!
- 3) LDR SP,[Rn,#imm]

4) LDR SP,[Rn]

5) LDR SP,[Rn,Rm]

The affected instructions that can result in the stack-pointer being loaded from an incorrect memory address are:

1) LDR SP,[Rn],#imm

2) LDR SP,[Rn,#imm]!

Conditions:

1) An LDR is executed, with SP/R13 as the destination.

2) The address for the LDR is successfully issued to the memory system.

3) An interrupt is taken before the data has been returned and written to the stack-pointer.

Implications:

Unless the load is being performed to Device or Strongly-Ordered memory, there should be no implications from the repetition of the load. In the unlikely event that the load is being performed to Device or Strongly-Ordered memory, the repeated read can result in the final stack-pointer value being different than had only a single load been performed.

Interruption of the two write-back forms of the instruction can result in both the base register value and final stack-pointer value being incorrect. This can result in apparent stack corruption and subsequent unintended modification of memory.

Workaround: Most compilers are not affected by this, so a workaround is not required.

However, for hand-written assembly code to manipulate the stack, both issues may be worked around by replacing the direct load to the stack-pointer, with an intermediate load to a general-purpose register followed by a move to the stack-pointer.

If repeated reads are acceptable, then the base-update issue may be worked around by performing the stack pointer load without the base increment followed by a subsequent ADD or SUB instruction to perform the appropriate update to the base register.

### e9004: Core: ITM can deadlock when global timestamping is enabled

#### Description: ARM ERRATA 806422

The Cortex-M4 processor contains an optional Instrumentation Trace Macrocell (ITM). This can be used to generate trace data under software control, and is also used with the Data Watchpoint and Trace (DWT) module which generates event driven trace. The processor supports global timestamping. This allows count values from a system-wide counter to be included in the trace stream.

When connected directly to a CoreSight funnel (or other component which holds ATREADY low in the idle state), the ITM will stop presenting trace data to the ATB bus after generating a timestamp packet. In this condition, the ITM\_TCR.BUSY register will indicate BUSY.

Once this condition occurs, a reset of the Cortex-M4 is necessary before new trace data can be generated by the ITM.

Timestamp packets which require a 5 byte GTS1 packet, or a GTS2 packet do not trigger this erratum. This generally only applies to the first timestamp which is generated.

Devices which use the Cortex-M optimized TPIU (CoreSight ID register values 0x923 and 0x9A1) are not affected by this erratum.

Workaround: There is no software workaround for this erratum. If the device being used is susceptible to this erratum, you must not enable global timestamping.

# e9005: Core: Store immediate overlapping exception return operation might vector to incorrect interrupt

**Description:** Arm Errata 838869: Store immediate overlapping exception return operation might vector to incorrect interrupt

Affects: Cortex-M4, Cortex-M4F

Fault Type: Programmer Category B Rare

Fault Status: Present in: r0p0, r0p1 Open.

The Cortex-M4 includes a write buffer that permits execution to continue while a store is waiting on the bus. Under specific timing conditions, during an exception return while this buffer is still in use by a store instruction, a late change in selection of the next interrupt to be taken might result in there being a mismatch between the interrupt acknowledged by the interrupt controller and the vector fetched by the processor.

**Configurations Affected** 

This erratum only affects systems where writeable memory locations can exhibit more than one wait state.

**Workaround:** For software not using the memory protection unit, this erratum can be worked around by setting DISDEFWBUF in the Auxiliary Control Register.

In all other cases, the erratum can be avoided by ensuring a DSB occurs between the store and the BX instruction. For exception handlers written in C, this can be achieved by inserting the appropriate set of intrinsics or inline assembly just before the end of the interrupt function, for example:

ARMCC:

•••

```
___schedule_barrier();
```

\_\_asm{DSB};

\_\_schedule\_barrier();

}

GCC:

\_asm volatile ("dsb 0xf" ::: "memory");

}

# e6940: Core: VDIV or VSQRT instructions might not complete correctly when very short ISRs are used

**Description:** Arm Errata 709718: VDIV or VSQRT instructions might not complete correctly when very short ISRs are used

Affects: Cortex-M4F

Fault Type: Programmer Category B

Fault Status: Present in: r0p0, r0p1 Open.

On Cortex-M4 with FPU, the VDIV and VSQRT instructions take 14 cycles to execute. When an interrupt is taken a VDIV or VSQRT instruction is not terminated, and completes its execution while the interrupt stacking occurs. If lazy context save of floating point state is enabled then the automatic stacking of the floating point context does not occur until a floating point instruction is executed inside the interrupt service routine.

Lazy context save is enabled by default. When it is enabled, the minimum time for the first instruction in the interrupt service routine to start executing is 12 cycles. In certain timing conditions, and if there is only one or two instructions inside the interrupt service routine, then the VDIV or VSQRT instruction might not write its result to the register bank or to the FPSCR.

**Workaround:** A workaround is only required if the floating point unit is present and enabled. A workaround is not required if the memory system inserts one or more wait states to every stack transaction.

There are two workarounds:

1) Disable lazy context save of floating point state by clearing LSPEN to 0 (bit 30 of the FPCCR at address 0xE000EF34).

2) Ensure that every interrupt service routine contains more than 2 instructions in addition to the exception return instruction.

### e11543: FlexCAN: Nominal Phase SJW incorrectly applied at CRC Delimiter

**Description:** During the reception of a CAN-FD frame when the Bit Rate Switch (BRS) is enabled, the Synchronization Jump Width (SJW) for the CRC Delimiter bit is incorrectly defined by the Nominal Phase SJW. The CAN specification stipulates that the CRC Delimiter bit should have a SJW set by the Data Phase SJW.

When a resynchronization event is triggered for the CRC delimiter bit (recessive in correct operation), the sample point will be adjusted by an amount as defined by the Nominal Phase SJW rather than the specified Data Phase SJW. This may result in the incorrect detection of a dominant bit leading to a CAN error frame. However, as the CRC delimiter bit position will only apply the SJW upon the detection of an unexpected dominant bit on the CAN bus, an error frame is already likely. For the case the SJW is applied at the CRC delimiter and a recessive bit is not detected, the receiving node will issue an error frame.

The CAN protocol is designed to handle resynchronization errors and hence the CAN bus will recover from the insertion of the incorrect SJW at the CRC delimiter. Upon detecting the error frame the transmitting node will re-transmit the frame.

The following FlexCAN configurations are not affected:

- Classical CAN frames (CAN 2.0B)
- CAN FD frames with bit rate switch disabled (BRS = 0)
- CAN FD frames with Nominal Phase SJW equal to Data Phase SJW
- CAN FD transmissions

Configuration for the FlexCAN:

- Nominal Phase SJW is configured by the Resync Jump Width bit in the CAN Control Register 1 (CAN\_CTRL1[RJW]) or by the Extended Resync Jump Width bit in the CAN Bit Timing Register (CAN\_CBT[ERJW])
- Data Phase SJW is configured by the Fast Resync Jump Width bit in the CAN FD Bit Timing Register (CAN\_FDCBT[FRJW])

**Workaround:** The robustness of the CAN protocol ensures that the receiver automatically recovers from the application of the incorrect SJW. The CAN protocol is designed to recover from resynchronization errors and hence any frame that is not correctly received will be re-sent by the transmitting node.

### e10856: FTM: Safe state is not removed from channel outputs after fault condition ends if SWOCTRL is being used to control the pin

- **Description:** If an FTM channel output is being controlled using the software output control register (FTM\_SWOCTRL) and fault detection is also enabled for the channel, then when a fault is detected the output is forced to its safe value. However, when the fault condition has been cleared, the channel output will stay in the safe state instead of reverting to the value programmed by the FTM\_SWOCTRL register.
- **Workaround:** If fault control is enabled while the software output control register is also being used (FTM\_SWOCTRL), then the FTM should be configured as follows:

-- FTM\_MODE[FAULTM] configured for manual fault clearing (0b10)

-- For devices that include the FTM\_CONF[NUMTOF] field, it must be cleared to 0b00000 (TOF set for each counter overflow). For FTM versions that don't include the FTM\_CONF[NUMTOF] field this doesn't apply.

The procedure below must be used in the TOF interrupt handler when a fault is detected to ensure that the outputs return to the value configured by FTM\_SWOCTRL.

1. Check the value of FTM\_FMS[FAULTF].

-- If FTM\_FMS[FAULTF] = 1 (fault occurred or is occurring), then set a variable to indicate that a fault was detected and continue to step 2.

-- If FTM\_FMS[FAULTF] = 0 but the fault variable is set (fault is not active, but was previously detected), continue to step 6.

2. Write the FTM\_OUTMASK register to set the bit or bits corresponding to any channels that are controlled by FTM\_SWOCTRL to temporarily inactivate the channel output.

3. Clear fault conditions by reading the FTM\_FMS register and then writing FTM\_FMS with all zeroes.

4. Clear the FTM\_SC[TOF] bit by reading the FTM\_SC register, then writing a 0 to FTM\_SC[TOF].

5. Exit the interrupt handler to skip following steps (they will execute the next time the TOF handler is called).

6. Clear the FTM\_SWOCTRL by writing all zeroes to it.

7. Write FTM\_SWOCTRL with the desired value again.

8. Clear the FTM\_OUTMASK bits that were set in step 2.

9. Clear the fault variable that was set in step 1 when the fault condition was originally detected.

10. Clear the FTM\_SC[TOF] bit by reading the FTM\_SC register, then writing a 0 to FTM\_SC[TOF].

# e10752: LPI2C: Slave Busy Flag may incorrectly read as zero when 10-bit address mode enabled in slave mode.

- **Description:** When operating in slave mode with 10-bit addressing enabled and an address match is detected on the first address byte, the Slave Busy Flag, SSR[SBF], may incorrectly read as zero.
- **Workaround:** When using the LPI2C in slave mode when 10-bit addressing is enabled, and the SSR[SBF] bit is read as a one, then the flag is set. If it is read as a zero, it must be read a second time and this second read will be the correct state of the bit.

# e10792: LPI2C: Slave Transmit Data Flag may incorrectly read as one when TXCFG is zero.

**Description:** When SCFGR1[TXCFG] = 0, the slave transmit data ready flag can incorrectly assert for one cycle.

Workaround: Set SCFGR1[TXCFG] = 1.

### e11097: LPSPI: Command word not loaded correctly when TXMSK=1

- **Description:** When the Transmit Command Register is written with TCR[TXMSK]=1 and the next write to the TX FIFO is another command, then the first command may not load correctly.
- **Workaround:** When writing the Transmit Command Register with TCR[TXMSK]=1, wait for the TX FIFO to go empty (FSR[TXCOUNT] = 0) before writing another command to the Transmit Command Register.

# e11089: LPSPI: In Continuous transfer mode with CPHA =1, WCF bit is not set for every word.

- **Description:** When Transmit Command Register is written with TCR[CONT]=1 and TCR[CPHA]=1, SR[WCF] bit flag is not set after data is transferred. Therefore polling for SR[WCF] flag to identify if data has been sent can cause MCU to be stuck.
- Workaround: When using continuous transfer mode TCR[CONT]=1 and TCR[CPHA]=1, do not use SR[WCF] flag to determine if data has been sent, fill up instead transmit FIFO with the following data without waiting for SR[WCF] flag to be set.

### e10655: LPSPI: Module Busy Flag may incorrectly read as zero in Master mode

**Description:** When operating in master mode, the Module Busy Flag, SR[MBF], may incorrectly read as zero. This applies only to master mode and not to slave mode.

Workaround: When using the LPSPI in master mode and the SR[MBF] bit is read as a one then the flag is set. If it is read as a zero, it must be read a second time and this second read will be the correct state of the bit.

# e10527: LPUART: Setting and immediately clearing SBK bit can result in transmission of two break characters

- **Description:** When the LPUART transmitter is idle (LPUART\_STAT[TC]=1), two break characters may be sent when using LPUART\_CTRL[SBK] to send one break character. Even when LUART\_CTRL[SBK] is set to 1 and cleared (set to 0) immediately.
- **Workaround:** To queue a single break character via the transmit FIFO, set LPUART\_DATA[FRETSC]=1 with data bits LPUART\_DATA[T9:T0]=0.

### e10716: RTC: Timer Alarm Flag can assert erroneously

- **Description:** Writing to the Time Alarm Register (RTC\_TAR) at the same time the RTC Seconds Register is incrementing can assert the Time Alarm Flag (RTC\_SR[TAF]) bit in the RTC Status Register.
- **Workaround:** Write the Time Alarm Register (RTC\_TAR) when the RTC Seconds Register is not incrementing. This can be when Time Counter Enable (RTC\_SR[TCE]) bit in the RTC Status Register is clear or within the RTC\_SR[TAF] interrupt routine.

Alternatively, if the RTC\_SR[TAF] is asserted following a write to the RTC\_TAR, then write the RTC\_TAR again.

### e10777: SCG: Corrupted status when the system clock is switching.

- **Description:** The SCG\_RCCR[SCS] and SCG\_HCCR[SCS] may have a corrupted status during the interval when the system clock is switching
- Workaround: The SCS field should be read twice by the software to ensure the system clock switch has completed.

# e11063: SMC: An asynchronous wakeup event during VLPS mode entry may result in possible system hang scenario.

- **Description:** When the bus clock is same system clock and an asynchronous wakeup occurs during a mode transition from RUN to VLPS or VLPR to VLPS, the MCU may hang in an undetermined state, which can only be recovered by a power-on reset event or a watchdog reset.
- **Workaround:** Before executing the transition to VLPS ensure that the PREDIV\_SYS\_CLK frequency / BUS\_CLK frequency configuration for RUN/VLPR mode is greater than or equal to 2.

For example: Assuming a PREDIV\_SYS\_CLK of 8 MHz and SCG\_RCCR[DIVCORE] = 0b0001 (divider of 2) and SCG\_RCCR[DIVBUS] = 0b0001 (divider of 1), (PREDIV\_SYS\_CLK = 8 MHz) / (BUS\_CLK = 4 MHz), a ratio of 1:2.

### e11114: SMC: invalid data might be fetched while accessing Flash in VLP modes

**Description:** VLPR and VLPS Low power modes are documented to work at System Clock and Core Clock at 4 Mhz and the Bus Clock at 4 MHz and DMA enabled from or to Flash memory. However any simultaneous access from any master (Core or DMA) to Dflash and Pflash may get invalid data while being in VLP modes and System clock, Core Clock and Bus Clock are above 1 Mhz

#### Workaround: There are two workarounds:

1. Restrict software to use either only Pflash or only Dflash only at a time in VLP modes for all masters (CPU,DMA). When switching from Pflash only access to Dflash only access let current DMA transactions accessing flash to complete and jump to SRAM location, wait for 40 cycles for the ongoing accesses to complete on the current flash before accessing dflash.

When switching from dflash only accesses to pflash only accesses let the current DMA transactions accessing dflash to complete

and wait for 40 cycles for accesses to complete on the dflash before accessing the pflash.

2. If both Pflash and Dflash needs to be accessed simultaneously, the VLP modes must be run with System Clock, Core Clock and Bus Clock of 1 MHz.

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